22222222222 222222222222	DDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDD		DDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDD	
CCC CCC CCC	DDD DDD DDD	DDD DDD DDD	DDD DDD DDD	DDD DDD DDD
CCC CCC	DDD DDD DDD	DDD DDD	DDD DDD DDD	DDD DDD DDD
CCC CCC CCC CCC	DDD DDD DDD	DDD DDD	DDD DDD DDD	DDD DDD
222	DDD DDD DDD DDD	DDD DDD DDD	DDD DDD DDD	DDD DDD DDD
000 000 000 000 000	DDD DDD DDDDDDDDDDDDD	DDD	DDD	DDD
2222222222	DDDDDDDDDDDD DDDDDDDDDDDD		DDDDDDDDDDDD DDDDDDDDDDDD DDDDDDDDDDDD	

22222222 22222222 22222222 22222222 2222	DDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDD	DDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDDD		BBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBB
BBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBB	3333333 3333333 3333333 3333333 3333333	2222222 22 22 22 22 22 22 22 22 22 22 2		

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TITLE: CDDLIB

CDD Bliss Library

VERSION: 'V04-000'

FACILITY: Common Data Dictionary

ABSTRACT:

This module is the library file used for compiling all other modules in the CDD facility.

**ENVIRONMENT:** 

AUTHOR: Jeff East, 22-Jan-80

MODIFIED BY:

7-May-81 (JAE) Added \$10\_SYNC macro.

XTITLE 'CDD Bliss Library'

```
16-SEP-1984 16:58:51.80 Page 2
CDDLIB.B32:1
    %SBTTL
                  'STRUCTURE DEFINITIONS'
         STRUCTURE DEFINITIONS
         On-disk Pointer Structure
                                                      DPTR$V_LINE
DPTR$V_PAGE
                       Logical Page
                                       ! Line !
                          number
                                       inumber!
                          +3
LITERAL
    DPTR$S_BLOCK_LENGTH = 3;
MACRO

$DPTR = BLOCK[DPTR$S_BLOCK_LENGTH,BYTE] FIELD (DPTR$Z_FIELDS)
FIELD DPTR$Z_FIELDS =
        DPTR$V_VALUE
DPTR$V_LINE
DPTR$V_PAGE
                                   = [0, 0, 24, 0],
= [0, 0, 7, 0],
= [0, 7, 17, 0]
    TES:
```

In-core Disk Pointer Structure

```
Internal: DKEY$V_LINE
file | Logical Page | Line | DKEY$V_PAGE
| number | number | DKEY$B_FILE
```

When a disk pointer is being passed around routines, it is passed as a Disk Key (DKEY), rather than just a disk pointer.

Disk Keys include all the information of a disk pointer, but also include a pointer to their file's FCB.

LITERAL DKEY\$S\_BLOCK\_LENGTH = 4;

SDKEY = BLOCK[DKEY\$S\_BLOCK\_LENGTH,BYTE] FIELD (DKEY\$Z\_FIELDS)

Line Index Format

31

Offset to blk

LINESW\_BLOCK\_OFFSET

Each block on a dictionary page is pointed to by a line index on that page. Disk pointers (DPTR) between the various blocks actually point to the block's line index. The line index is then used to locate the physical dictionary block.

The PAGE\$INDEX macro allows the programmer to access a line index while it resides on a dictionary page. It also uses the LINE\$Z\_FIELDS to access the individual fields within a line index.

LINE\$S\_BLOCK\_LENGTH = 2;

MACRO
\$LINE = BLOCK[LINE\$S\_BLOCK\_LENGTH,BYTE] FIELD (LINE\$Z\_FIELDS)

FIELD LINE\$Z\_FIELDS =

SET

LINE\$W\_BLOCK\_OFFSET = [0, 0, 16, 0]

TES;

Page Format

31	15	0	+ +0	
Page Checksum				PAGESL_CHECKSUM
Page's cluster seq #	Page Nu	Page Number		PAGESV_NUMBER PAGESB_CLUSTER_IO_SEQ
Size of next Next Related page group  Page's Prior Related Pag seq #			+8	PAGE\$L_NEXT_PAGE PAGE\$V_NEXT_PAGE PAGE\$V_NEXT_SIZE PAGE\$V_PRIOR_PAGE PAGE\$B_GROUP_IO_SEQ
			+16	
Number of Free space line indices		space	+20	PAGESW_FREE_SPACE PAGESW_INDICES PAGESV_DATA
	Data area /\/\/\/\/ ne index area		•	
			+512	PAGESV_INDEX_BASE

Each page in the dictionary file starts with a page header. The next and prior related page pointers are used for linking page groups. The following types of page groups exist:

1) Lockable page group.
Each named entity (and history list) owns exactly one lockable page group. These groups contain all unamed offspring of the name entiry (or history list).

A page group may only be accessed through its portal page. A group's portal page is the page on which the named entity (or history list head) resides. If a group's portal page is locked, then no pages in the group may be accessed.

The PAGE\$INDEX structure is used to access a line index entry on a page.

A page whose checksum is zero is a locked page, and indicates that

the sub-tree below it is incomplete and in a transient state. Such pages may never be read.

Each page has two page sequence numbers. The cluster sequence number must be the same for all pages in the cluster. The group sequence number must be the same for all pages in an I/O group.

The cluster I/O sequence number is bumped whenever more than one group in the cluster is written. The group I/O sequence number is bumped whenever the group is written. These sequence numbers enable the detection of incomplete cluster unstage operations.

```
LITERAL
           PAGESS_BLOCK_LENGTH = 512:
           $PAGE = BLOCK[PAGE$S_BLOCK_LENGTH,BYTE] FIELD (PAGE$Z_FIELDS)
FIELD PAGESZ_FIELDS =
           SET
                    PAGESL_CHECKSUM
PAGESV_NUMBER
PAGESB_CLUSTER_IO_SEQ
PAGESL_NEXT_PAGE
PAGESV_NEXT_PAGE
PAGESV_NEXT_SIZE
PAGESV_PRIOR_PAGE
PAGESW_FREE_SPACE
PAGESW_INDICES
PAGESV_DATA
PAGESV_INDEX_BASE
                                                                                     = [0, 0, 32, 0],

= [4, 0, 17, 0],

= [7, 0, 8, 0],

= [8, 0, 32, 0],

= [8, 0, 17, 0],

= [8, 17, 15, 0],

= [12, 0, 17, 0],

= [15, 0, 8, 0],

= [16, 0, 16, 0],

= [18, 0, 16, 0],

= [20, 0, 0, 0],

= [512, 0, 0, 0]
           TES:
LITERAL
          PAGESK_BASE = BLOCK[O, PAGESV_DATA;

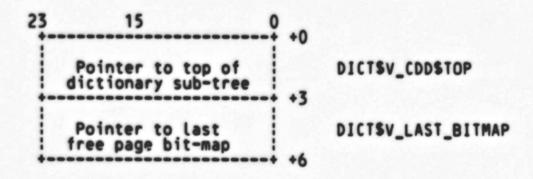
PAGESS_BLOCK_LENGTH, BYTE],

PAGESS_INDEX_BASE = PAGESS_BLOCK_LENGTH - PAGESK_BASE,

PAGESS_INDEX_BASE = BLOCK[O, PAGESV_INDEX_BASE;

PAGESS_BLOCK_LENGTH, BYTE];
STRUCTURE
          PAGESINDEX[I, O, P, S, E] = (PAGESINDEX+PAGESS_INDEX_BASE+O-((I)*LINE$S_BLOCK_LENGTH)) <P,S,E>;
```

Dictionary Header Block



Each dictionary uses page T as its dictionary header page. The dictionary header block immediately follows the page header on page 1.

DICT\$V\_CDD\$TOP points to the root of the tree/sub-tree contained in the dictionary file.

DICTSV\_LAST\_BITMAP points to the last free page bit-map entry.

DICTSS\_BLOCK\_LENGTH = 6 + PAGESK\_BASE;

SDICT = BLOCK[DICTSS\_BLOCK\_LENGTH, BYTE] FIELD (PAGESZ\_FIELDS)

FIELD DICTSZ\_FIELDS =

DICTSV\_CDDSTOP = [0+PAGESK\_BASE, 0, 24, 0],
DICTSV\_LAST\_BITMAP = [3+PAGESK\_BASE, 0, 24, 0]

! \*

Block types

The ordering and clustering of these block types is significant.

Ail block types must be contiguous without any holes and bounded by the symbols BLK\$K\_TYPE\_FIRST and BLK\$K\_TYPE\_LAST.

All node and NAME block types must be contiguous and bounded by the symbols BLK\$K\_TYPE\_NODE\_FIRST and BLK\$K\_TYPE\_NODE\_LAST.

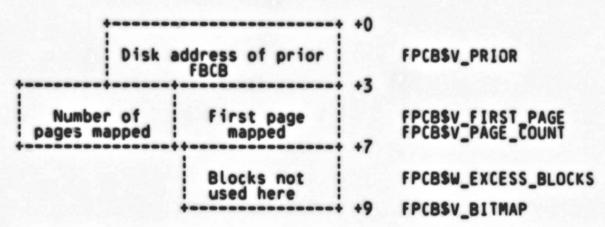
To add more node or NAME block types, insert them at the front of the table.

To add any other block type, append them to the end of the table.

CHANGING THE VALUES OF ANY OF THE ACTUAL BLOCK TYPES WILL INVALIDATE EXISTING DICTIONARIES.

```
LITERAL
        BLKSK_TYPE_FIRST
                                                                                            = 101, ! Lowest block type
                                                                                            = 101, ! First node or NAM block
                  BLK$K_TYPE_NODE_FIRST
                          BLKSK_TYPE_DIR_NAM
BLKSK_TYPE_FIL_NAM
BLKSK_TYPE_TERM_NAM
BLKSK_TYPE_DIR_NODE
BLKSK_TYPE_TERM_NODE
                                                                                            = 101.
= 102.
= 103.
= 104.
= 105.
                                                                                                                   Directory NAM block
                                                                                                                    File NAM block
                                                                                                                   Terminal NAM block
                                                                                                                   Directory node block
                                                                                                               ! Terminal node block
                  BLK$K_TYPE_NODE_LAST
                                                                                            = 105, ! Last node or NAM block
       BLKSK TYPE BITMAP
BLKSK TYPE ENTITY ATT
BLKSK TYPE ENTITY LIST
BLKSK TYPE ENTITY LIST
BLKSK TYPE NULL ATT
BLKSK TYPE NUM ATT
BLKSK TYPE STRING ATT
BLKSK TYPE STRING LIST
BLKSK TYPE STRING LIST
BLKSK TYPE STRING LIST
BLKSK TYPE STRING SEG
BLKSK TYPE ACLC
BLKSK TYPE ACLC
BLKSK TYPE ACLC
                                                                                            = 106.
= 107.
                                                                                                                   free page bitmap
                                                                                                                   Entity attribute block
Entity list block
Entity list attribute block
                                                                                            = 108.
= 109.
                                                                                            = 109,
= 110,
= 111,
= 112,
= 113,
= 115,
= 116,
= 117,
= 118,
= 120,
= 120;
                                                                                                                   Null attribute block
Numeric attribute block
Short string attribute block
String attribute block
String list block
String list attribute block
                                                                                                                    String segment block
Text block
                                                                                                                   Access control list entry
Access control list criterion
                                                                                                                   Highest block type
```

free Page Control Block (FPCB) free Page Bit Map (FPBM)



free pages in the dictionary file are controlled by the free Page Control Blocks (FPCB) and the Free Page Bit Map (FPBM). Each free page in the file has its corresponding bit set on.

The disk's cluster size may make it impossible to use all of the allocated blocks. If so, the un-used blocks are tallied in FPCB\$W\_EXCESS\_BLOCKS, and will be used when the file is next extended.

The FPCBs are chained as the dictionary file grows in size.

```
LITERAL

FPCBSS_BLOCK_LENGTH = BLK$K_BASE + 9;

STRUCTURE

FPCB[O, P, S, E; BLOCKS] = [FPCBSS_BLOCK_LENGTH+(BLOCKS+7)/8]

(FPCBV_PCBSV_FIELDS =

SET

FPCBSV_PRIOR

FPCBSV_PRIOR

FPCBSV_PAGE TOUNT

FPCBSV_PAGE TOUNT

FPCBSW_EXCESS_BLOCKS

FPCBSV_BITMAP

TES;

LITERAL

FPCBSK_BASE = FPCBSS_BLOCK_LENGTH,
FPCBSS_BITMAP = FPCB[O, FPCBSV_BITMAP];

STRUCTURE

FPBM[I] = (FPBM+FPCBSS_BITMAP)<I-1, 1, 0>;
```

```
16-SEP-1984 16:58:51.80 Page 12
CDDLIB.B32:1
          Name Block
           31
                               15
                        Pointer to next NAME
                                                          NAME$V_SIBLING
                               (sibling)
                       Pointer to name block
                                                          NAMESV_NAME
                        Pointer to node block or forwarding
                                                          NAMESV_NODE or NAMESV_FILE
                               file name
          There are three types of NAM Blocks, Directory, File, and Terminal.
          Directory NAM Block is made up of a Block Header + NAM Block.
          File NAM Block is made up of a Block Header + NAM Block.
          Terminal NAM Block is made up of a Block Header + NAM Block + Terminal
         NAM Block.
NAMESS_BLOCK_LENGTH = 9 + BLKSK_BASE;
MACRO
    SNAME = BLOCK[NAMESS_BLOCK_LENGTH, BYTE] FIELD (BLK$Z_FIELDS, NAME$Z_FIELDS)
     %;
FIELD NAMESZ_FIELDS =
         NAMESV_SIBLING
NAMESV_NAME
NAMESV_NODE
NAMESV_FILE
                                      = [0+BLK$K_BASE, 0, 24,
= [3+BLK$K_BASE, 0, 24,
= [6+BLK$K_BASE, 0, 24,
= [6+BLK$K_BASE, 0, 24,
     TES.
LITERAL
```

= NAME\$S\_BLOCK\_LENGTH;

NAMESK\_BASE

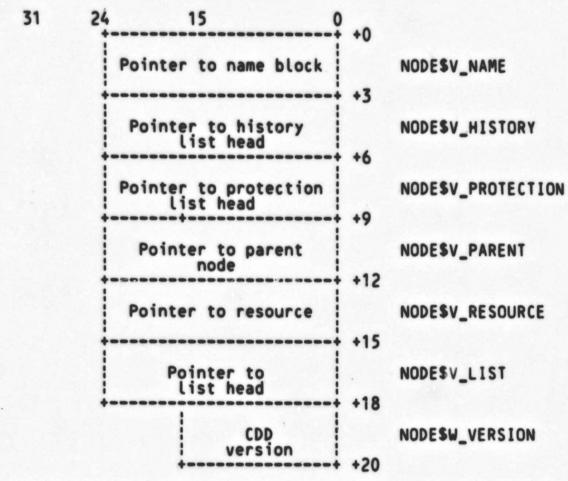
```
16-SEP-1984 16:58:51.80 Page 13
 CDDLIB.B32;1
          Node Nam Block
           31
                                                       NNAM$V_RESOURCE
                       Pointer to resource
         Node NAM Block is made up of a Block Header + NAM Block + Node NAM Block.
LITERAL NNAMSS_BLOCK_LENGTH = 3 + NAMESK_BASE;
MACRO

$NNAM = BLOCK[NNAM$S_BLOCK_LENGTH,BYTE] FIELD (BLK$Z_FIELDS,

NAME$Z_FIELDS)

NNAM$Z_FIELDS)
     X;
FIELD NNAMSZ_FIELDS =
         NNAM$V_RESOURCE = [0+NAME$K_BASE, 0, 24, 0]
     TES:
LITERAL
NNAMSK_BASE
                           = NNAM$S_BLOCK_LENGTH;
```

Node Block



There are two types of Node Blocks Directory and Terminal.

Directory Node Block is made up of a Block Header + Node Block

Terminal Node Block is made up of a Block Header + Node Block

```
LITERAL
NODE$S_BLOCK_LENGTH = 20+ BLK$K_BASE;

MACRO
$NODE = BLOCK[NODE$S_BLOCK_LENGTH,BYTE] FIELD (BLK$Z_FIELDS,
NODE$Z_FIELDS)

X;

FIELD NODE$Z_FIELDS =

SET

NODE$V_ORDERED
NODE$V_NAME
NODE$V_NAME
NODE$V_HISTORY

= [0+BLK$K_BASE. 0. 24. 0].
```

```
CDDLIB.B32;1

NODE$V_PROTECTION = [6+BLK$K_BASE, 0, 24, 0],
NODE$V_PARENT = [9+BLK$K_BASE, 0, 24, 0],
NODE$V_PARENT = [12+BLK$K_BASE, 0, 24, 0],
NODE$V_LIST = [15+BLK$K_BASE, 0, 24, 0],
NODE$W_VERSION = [18+BLK$K_BASE, 0, 16, 0],
TES;

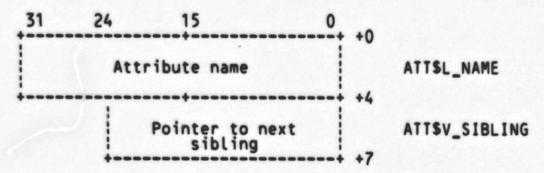
LITERAL NODE$K_BASE = NODE$S_BLOCK_LENGTH;

The following flag occurs in the BLK$B_FLAGS field of a directory node.

LITERAL NODE$M_ORDERED = 1^1 - 1^0; ! List elements are sorted by name
```

!+

Common Attribute Block



Each attribute has the common attribute block immediately following the universal block header.

There are six types of Attribute Blocks Entity Attribute Block, List Attribute Block, Numeric Attribute Block, Null Attribute Block, Short String Attribute Block, and String Attribute Block.

Entity Attribute Block is made up of a Block Header + Attribute Block + Entity Attribute Block.

List Attribute Block is made up of a Block Header + Attribute Block + List Attribute Block.

Numeric Attribute Block is made up of a Block Header + Attribute Block + Numeric Attribute Block.

Null Attribute Block is made up of a Block Header + Attribute Block.

Short String Attribute Block is made up of a Block Header + Attribute Block + Short String Attribute.

String Attribute Block is made up of a Block Header + Attribute Block + String Attribute Block.

```
LITERAL

ATT$S_BLOCK_LENGTH = 7 + BLK$K_BASE;

MACRO

$ATT = BLOCK[ATT$S_BLOCK_LENGTH,BYTE] FIELD (ATT$Z_FIELDS, BLK$Z_FIELDS)

%;

FIELD ATT$Z_FIELDS =

SET

ATT$L_NAME
ATT$V_SIBLING

TES;

LITERAL
```

CDDL18.832;1

16-SEP-1984 16:58:51.80 Page 17

ATTSK\_BASE

= ATT\$S\_BLOCK\_LENGTH;

```
16-SEP-1984 16:58:51.80 Page 18
 CDDLIB.B32;1
          Entity Attribute Block
           31
                             15
                        Pointer to first attribute
                                                       ENTSV_FIRST_ATT
         Entity Attribute Block is made up of a Block Header + Attribute Block + Entity Attribute Block.
LITERAL ENTSS_BLOCK_LENGTH = 3 + ATTSK_BASE;
MACRO
SENT = BLOCK[ENT$S_BLOCK_LENGTH,BYTE] FIELD (ATT$Z_FIELDS,
BLK$Z_FIELDS)
     %;
FIELD ENT$Z_FIELDS =
    TES; ENTSV_FIRST_ATT
                                    = [0+ATT$K_BASE, 0, 24, 0]
LITERAL
ENT$K_BASE
                           = ENT$S_BLOCK_LENGTH;
```

```
16-SEP-1984 16:58:51.80 Page 19
CDDLIB.B32:1
          List Attribute Block
           31
                              15
                          Pointer to first
                                                          LIST$V_FIRST_LST
                            lst segment
                                 Total number of cells
                                                          LIST$W_CELL_COUNT
         List Attribute Block is made up of a Block Header + Attribute Block + List Attribute Block.
LIST$S_BLOCK_LENGTH = 5 + ATT$K_BASE;
MACRO

$LIST = BLOCK[LIST$S_BLOCK_LENGTH,BYTE] FIELD (ATT$Z_FIELDS,

BLK$Z_FIELDS)
     %;
FIELD LIST$Z_FIELDS =
                                       = [0+ATT$K_BASE, 0, 24, 0],
= [3+ATT$K_BASE, 0, 16, 0]
         LISTSV_FIRST_LST
LISTSW_CELL_COUNT
     TES:
LITERAL
LISTSK_BASE
                             = LIST$S_BLOCK_LENGTH;
```

```
16-SEP-1984 16:58:51.80 Page 20
CDDLIB.B32;1
         Numeric Attribute Block
                   24
                             15
                                                      NAT$L_VALUE
                Numeric Attribute Value
         Numeric Attribute Block is made up of a Block Header + Attribute Block + Numeric Attribute Block.
         Numeric attribute block contains the
         value of the numeric attribute.
LITERAL NATSS_BLOCK_LENGTH = 4 + ATTSK_BASE;
MACRO
    SNAT = BLOCK[NATSS_BLOCK_LENGTH, BYTE] FIELD (ATTSZ_FIELDS, BLK$Z_FIELDS, NATSZ_FIELDS)
    %:
FIELD NATSZ_FIELDS =
         NAT$L_VALUE
                                    = [0+ATT$K_BASE,0,32,0]
    TES:
LITERAL
NATSK_BASE
                                    = NAT$S_BLOCK_LENGTH;
```

```
16-SEP-1984 16:58:51.80 Page 21
CDDLIB.B32:1
            Short String Attribute Block
            31
                                                                       SSAST_STRING
                                                   String
            Short String Attribute Block is made up of a Block Header + Attribute Block + Short String Attribute.
           Strings whose total length is between 0 and 255 bytes are usually stored using the short string attribute block. If a string is too long to be stored in an SSA, then it is stored using a normal string attribute block (STR).
LITERAL
     SSASS_BLOCK_LENGTH = 0 + ATTSK_BASE;
     $SSA = BLOCK[SSA$S_BLOCK_LENGTH, BYTE] FIELD (ATT$Z_FIELDS, BLK$Z_FIELDS, SSA$Z_FIELDS)
      %;
FIELD SSASZ_FIELDS =
           SSAST_STRING
                                               = [0+ATT$K_BASE, 0, 0, 0]
      TES:
LITERAL
                                   = SSA$S_BLOCK_LENGTH;
      SSASK_BASE
```

```
16-SEP-1984 16:58:51.80 Page 22
CDDLIB.B32:1
          String Attribute Block
            31
                     24
                                15
                           Pointer to first string segment
                                                             STR$V_STRING
                                   Total string
                                                             STR$W_LENGTH
                                      length
          String Attribute Block is made up of a Block Header + Attribute Block + String Attribute Block.
          The string is broken into one or more segments. The string segments are stored in SEG blocks.
STRSS_BLOCK_LENGTH = 5 + ATTSK_BASE;
MACRO

$STR = BLOCK[STR$S_BLOCK_LENGTH,BYTE] FIELD (ATT$Z_FIELDS,

BLK$Z_FIELDS)
     %;
FIELD STR$Z_FIELDS =
          STRSV_STRING
STRSW_LENGTH
                                        = [0+ATT$K_BASE, 0, 24, 0].
= [3+ATT$K_BASE, 0, 16, 0]
     TES:
LITERAL
     STR$K_BASE
                              = STR$S_BLOCK_LENGTH;
```

```
16-SEP-1984 16:58:51.80 Page 23
CDDLIB.B32:1
           LST Segment Block
            31
                                  15
                                                                 LST$V_NEXT_LST
                           Pointer to next lst
                                                                 LST$W_CELL_COUNT
                                   Number of cells:
                                           here
                                                                LSTST_CELLS
           There are two types of LST Segment Blocks Entity List Block and String List Attribute Block.
           Entity List Block is made up of a Block Header + LST Block + Entity List Block.
           String List Attribute Block is made up of a Block Header + LST Block + String List Attribute Block.
LITERAL
LST$S_BLOCK_LENGTH = 5 + BLK$K_BASE;
MACRO
SLST = BLOCK[LST$S_BLOCK_LENGTH,BYTE] FIELD (BLK$Z_FIELDS),
LST$Z_FIELDS)
     %;
FIELD LST$Z_FIELDS =
          LSTSV_NEXT_LST
LSTSW_CELL_COUNT
LSTST_CELLS
                                           = [0 + BLK$K_BASE, 0, 24, 0],
= [3 + BLK$K_BASE, 0, 16, 0],
= [5 + BLK$K_BASE, 0, 0, 0]
      TES:
LITERAL
LSTSK_BASE
                                = LST$S_BLOCK_LENGTH;
```

```
16-SEP-1984 16:58:51.80 Page 24
CDDLIB.B32:1
         Entity List Block
          31
                             15
                       Pointer to first attribute
                                                     ELST$V_ATT
         Entity List Block is made up of a Block Header + LST Block + Entity List Block.
LITERAL
ELSTSS_BLOCK_LENGTH = 3;
MACRO

SELST = BLOCK[ELST$S_BLOCK_LENGTH,BYTE] FIELD (ELST$Z_FIELDS)

%;
FIELD ELST$Z_FIELDS =
        ELST$V_ATT
                          = [0, 0, 24, 0]
    TES;
LITERAL
ELSTSK_BASE
                       = ELST$S_BLOCK_LENGTH;
STRUCTURE
    ELSTSELM[]] = (ELSTSELM + LSTSK_BASE + (I * ELSTSS_BLOCK_LENGTH));
```

```
16-SEP-1984 16:58:51.80 Page 25
CDDLIB.B32:1
          String List Attribute Block
           31
                    24
                               15
                                          ----+ +0
                                                          SLST$V_STRING
                          Pointer to first
                           string segment
or text block
                                                          SLST$W_LENGTH
                                 Total string
                                     length
                                                           SLST$T_STRING
          String List Attribute Block is made up of a Block Header + LST Block + String List Attribute Block.
LITERAL
SLST$S_BLOCK_LENGTH = 5;
MACRO
     $SLST = BLOCK[SLST$S_BLOCK_LENGTH,BYTE] FIELD (SLST$Z_FIELDS)
FIELD SLST$Z_FIELDS =
          SLST$V_STRING
SLST$W_LENGTH
SLST$T_STRING
                                       = [0, 0, 24, 0],
= [3, 0, 16, 0],
= [5, 0, 0, 0]
     TES:
LITERAL
     SLST$K_BASE
                             = SLST$S_BLOCK_LENGTH;
STRUCTURE
     SLSTSELM[1] =
          (SLSTSELM + LSTSK_BASE + (I * SLST$S_BLOCK_LENGTH));
```

```
16-SEP-1984 16:58:51.80 Page 26
CDDLIB.B32:1
            String Segment Block
             31
                                                          0
                                Pointer to next
                                                                      SEG$V_NEXT
                                string segment
                                                                     SEGST_STRING
            String Segment Block is made up of a Block Header + String Segment Block.
           Each string attribute points to zero or more string segment blocks. Each block contains a portion of the whole string. The final string is found by catenating all the string segments together.
LITERAL
     SEG$S_BLOCK_LENGTH = 3 + BLK$K_BASE;
MACRO
     $SEG = BLOCK[SEG$S_BLOCK_LENGTH,BYTE] FIELD (BLK$Z_FIELDS, SEG$Z_FIELDS)
      %;
FIELD SEG$Z_FIELDS =
           SEGSV_NEXT
SEGST_STRING
                                              = [0+BLK$K_BASE, 0, 24, 0],
= [3+BLK$K_BASE, 0, 0, 0]
LITERAL
SEG$K_BASE
                                   = SEG$S_BLOCK_LENGTH;
```

```
16-SEP-1984 16:58:51.80 Page 27
CDDLIB.B32:1
        Text Block
                              +----+ +0 TEXT$T_STRING
        31
                                 Text
       Text Block is made up of a Block Header + Text Block.
LITERAL TEXT$S_BLOCK_LENGTH = 0 + BLK$K_BASE;
MACRO
STEXT = BLOCK[TEXT$S_BLOCK_LENGTH,BYTE] FIELD (BLK$Z_FIELDS)
    %;
FIELD TEXT$Z_FIELDS =
   TEXTST_STRING = [0+BLK$K_BASE, 0, 0, 0]
LITERAL
TEXT$K_BASE
                  = TEXT$S_BLOCK_LENGTH;
```

```
Access Control List Entry (ACL)
 31
             Ptr to next ACL
                                           ACL$V_NEXT
      Rights to be granted
                                           ACL$L_GRANT
       Rights to be denied
                                           ACL$L_DENY
      Rights to be banished
                                           ACL$L_BANISH
   UIC group # : UIC member #
                                           ACL$W_UIC_MEMBER
                                           ACL$W_UIC_GROUP
         ! Ptr to first ACLC block!
                                           ACL$V_FIRST_ACLC
Each node has an Access Control List made up of zero or more Access Control List Entries (ACL).
The ACL block is appended to a BLK to make an access control
list entry.
```

CDDLIB.B32;1

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LITERAL ACLSK\_BASE

= ACL\$S\_BLOCK\_LENGTH;

```
Access Control List Criterion (ACLC)
```

```
Type Ptr to next ACLC ACLC$V_NEXT ACLC$B_TYPE +4 ACLC$T_STRING
```

Each ACL may have one or more Access Control List Criterion (ACLC) chained from it.

Each ACLC specifies one user identification criterion for the ACL entry. The user must match all the criteria for the ACL entry to apply to him.

An ACLC is appended to a BLK to form the criterion block.

```
ACLC$S_BLOCK_LENGTH = 4+BLK$K_BASE;
```

SACLC = BLOCK[ACLC\$S\_BLOCK\_LENGTH, BYTE] FIELD (BLK\$Z\_FIELDS, ACLC\$Z\_FIELDS)

FIELD ACLCSZ\_FIELDS =

ACLC\$V\_NEXT = [0+BLK\$K\_BASE, 0, 24, 0],
ACLC\$B\_TYPE = [3+BLK\$K\_BASE, 0, 8, 0],
ACLC\$T\_STRING = [4+BLK\$K\_BASE, 0, 0, 0]

LITERAL ACLCSK\_BASE

= ACLC\$S\_BLOCK\_LENGTH;

```
User Control Block (UCB)
```

A user's context pointer is an origin 1 index into the CCBs in the UCB list. Each UCB can point to UCB\$K\_CCB number of CCBs. If the user passes a context number that doesn't map to an active CCB, we tell him it's an invalid context pointer.

Except in some extreme case, any image will probably never have more than 4 users active at any one time. But we can handle more!

```
LITERAL

UCB$K_CCB_BASE

UCB$K_CCB_BASE

UCB$K_CCB_BASE

UCB$K_CCB_BASE

UCB$K_CCB_BASE

UCB$K_CCB_BASE + UCB$K_CCB * 4;

MACRO

$UCB = BLOCK[UCB$S_BLOCK_LENGTH,BYTE] FIELD (UCB$Z_FIELDS)

%;

FIELD

UCB$Z_FIELDS =

UCB$A_NEXT

UCB$W_FIRST

UCB$W_FIRST

UCB$W_LAST

UCB$V_CCB

UCB$V_CCB

UCB$V_CCB

UCB$C_FIRST = BLOCK[O, UCB$W_FIRST; UCB$S_BLOCK_LENGTH, BYTE];

STRUCTURE

UCB$CCB[I, O, P, S, E] = (UCB$CCB+UCB$K_CCB_BASE+

UCB$CCB[I, O, P, S, E] = (UCB$CCB+UCB$CCB_BASE+

UCB$CCB_BASE +

UC
```

Pool Header Block (PHB)

```
Number | Low | PHB$B_POOL_TYPE | Of | Slot | Pool | PHB$B_LOW_SLOT | PHB$B_SLOTS | PHB$B_SLOTS |

Ptr to last MCB in the pool | PHB$A_LAST_MCB | PHB$A_LAST_MCB |

Number of bytes in each | Secondary extent | PHB$L_EXTENT_SIZE | PHB$V_FREE_BLOCK_LIST
```

Pools are used to provide space for various in-core block types. Blocks that are related are usually allocated from the same pool. This is done because pools provide good locality of reference and this allocation scheme reduces page faults.

Each pool consists of one or more extents. Each extent consists of a Memory Control Block and the rest of the space allocated to the extent. Each pools has a Pool Header Block associated with it. This block identifies the MCB list, as well as contains the pool's free block list.

The free block list is a vector of linked lists. Each slot in the vector corresponds to a block type. When a block is freed, it is linked into its associated free list. When a block is requested, its free list is checked to see if it is non-empty. If so, then a block is allocated from it. Otherwise, a block is allocated from one of the pool's extents.

```
CDDLIB.B32:1
```

TES:

```
Memory Control Block (MCB)
```

```
Pointer to prior memory extent MCB$A_PRIOR

Remaining size Size of this of this extent extent HCB$W_ALLOC_SIZE MCB$W_FREE_SIZE
```

Each pool has one or more Memory Control Blocks associated with it. This list serves two purposes:

- 1) It points to each memory extent we asked LIB\$GET\_VM for.
- It keeps track of memory that has been allocated for the pool, but never used.

```
LITERAL MCBSS_BLOCK_LENGTH = 8;

MACRO SMCB = BLOCK[MCBSS_BLOCK_LENGTH,BYTE] FIELD (MCBSZ_FIELDS)
%;

FIELD MCBSZ_FIELDS =

MCBSA_PRIOR = [0, 0, 32, 0],

MCBSW_ALLOC_SIZE = [4, 0, 16, 0],

MCBSW_FREE_SIZE = [6, 0, 16, 0]
```

```
Hash Control Entry (HCE)
             31
                                    15
                    Forward homonym pointer
                                                                     HCESA_NEXT
                   Backwards homonym pointer
                                                                     HCESA_PRIOR
            # blocks ever
stored in bucket
                                                                     HCESW_HITS
HCESW_BLOCKS
                                       # times bucket!
                                          accessed
              Max blocks ever! # blocks now in bucket chain in bucket chain!
                                                                     HCESW_CUR_LENGTH
           HCEs are found in each user's CCB. Each HCE forms one hash bucket. We keep some statistics on the buckets in an attempt to tune the hash function.
LITERAL
     HCE$S_BLOCK_LENGTH = 16;
     SHCE = BLOCK[HCESS_BLOCK_LENGTH, BYTE] FIELD (HCESZ_FIELDS)
FIELD HCESZ_FIELDS =
           HCESA_NEXT
HCESA_PRIOR
HCESW_HITS
HCESW_BLOCKS
HCESW_CUR_LENGTH
HCESW_MAX_LENGTH
      TES:
LITERAL
     HCESK_HOMONYM_LIST
                                              = BLOCK[O, HCE$A_NEXT; ,BYTE];
```

```
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CDDLIB.B32:1
            Hash Control Block (HCB)
             31
                                    15
                    forward homonym pointer
                                                                     HCB$A_NEXT
                   Backwards homonym pointer
                                                                     HCB$A_PRIOR
                             Hash key
                                                                     HCB$V_HASH_KEY
                        Rest of hash key
                   Pointer to object block
                                                                     HCB$A_BLOCK
LITERAL
     HCB$K_PCB_NUMBER = 1,
HCB$K_LCCB_CODE = 2,
HCB$K_LCCB_ADDRESS = 3,
HCB$K_KEY_ENGTH = 8,
HCB$S_BLOCK_LENGTH = 20;
                                              ! Hash type is page number
! Hash type is location code
! Hash type is entity disk address
! Size of a hash key
     $HCB = BLOCK[HCB$S_BLOCK_LENGTH,BYTE] FIELD (HCB$Z_FIELDS)
MACRO
FIELD HCB$Z_FIELDS =
      SET
           HCBSA_NEXT
HCBSA_PRIOR
HCBSV_HASH_KEY
HCBSA_BLOCK
      TES:
LITERAL
     HCB$K_HOMONYM_LIST
                                             = BLOCK[O, HCB$A_NEXT; , BYTE];
```

# Context Control Block (CCB)

31	15	0	+ +0	
		Flags		CCB\$W_FLAGS
Poin in t	ter to first emporary lock	LCB list	+4	CCB\$A_FIRST_LOCK
Poi in t	nter to last emporary lock	LCB list		CCB\$A_LAST_LOCK
User		r ID for manager	+12	CCB\$W_LOCK_ID CCB\$W_USER_ID
Ne numb	Next location code number to be assigned			CCB\$L_NEXT_LCC
Ptr	Ptr to NCB for file's CDD\$TOP node		+20	CCB\$\$A_CDD\$TOP_NCB
NC	NCB of CDD\$LOGIN		+24	CCB\$A_LOGIN_NCB
Pt	Ptr to current FCB for user		+28	CCB\$A_CURRENT_FCB
Ptr	Ptr to highest cluster		+32	CCB\$A_FIRST_CIB
Pt	Ptr to last highest cluster in cache		+36	CCB\$A_LAST_CIB
St	Status of current transaction		+40	CCB\$L_STATUS
	Ptr to CCB's pool header		+44	CCB\$A_POOL
	Ptr to first			CCB\$A_FIRST_SPB

```
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 CDDLIB.B32:1
                           Security Preservation Block :
                                         Ptr to last
                                                                                                             CCB$A_LAST_SPB
                           Security Preservation Block
                                                                                                 +56 CCB$V_HASH_BASE
                   The hash table consists of a number of hash entries.
                  See the HCE description for the format of these entries.
LITERAL
         CCB$K_HASH_TABLE_SIZE = 151,
CCB$S_BLOCK_LENGTH = 56 + (CCB$K_HASH_TABLE_SIZE * HCE$S_BLOCK_LENGTH);
MACRO
         $CCB = BLOCK[CCB$S_BLOCK_LENGTH,BYTE] FIELD (CCB$Z_FIELDS)
FIELD_CCB$Z_FIELDS =
                 CCBSW_FLAGS
CCBSV_CORRUPT
CCBSA_FIRST_LOCK
CCBSW_LOCK_ID
CCBSW_USER_ID
CCBSL_NEXT_LCC
CCBSA_CDDSTOP_NCB
CCBSA_LOGIN_NCB
CCBSA_LOGIN_NCB
CCBSA_LOGIN_NCB
CCBSA_LOGIN_NCB
CCBSA_FIRST_CIB
CCBSA_LAST_CIB
CCBSA_LAST_CIB
CCBSA_LAST_CIB
CCBSA_FIRST_SPB
CCBSA_LAST_SPB
CCBSA_LAST_SPB
CCBSA_LAST_SPB
CCBSA_LAST_SPB
         SET
                                                                       = [0, 0, 16, 0],

= [0, 0, 1, 0],

= [4, 0, 32, 0],

= [8, 0, 32, 0],

= [12, 0, 16, 0],

= [14, 0, 16, 0],

= [16, 0, 32, 0],

= [20, 0, 32, 0],

= [24, 0, 32, 0],

= [28, 0, 32, 0],

= [32, 0, 32, 0],

= [40, 0, 32, 0],

= [44, 0, 32, 0],

= [44, 0, 32, 0],

= [44, 0, 32, 0],

= [48, 0, 32, 0],

= [52, 0, 32, 0],

= [52, 0, 32, 0],

= [52, 0, 32, 0],

= [56, 0, 0, 0]
                                                                                                                               ! Stream is corrupt
         TES:
LITERAL
         CCB$K_LOCK_LIST = BLOCK[O, CCB$A_FIRST_LOCK; , BYTE],
CCB$K_CLUSTER_LIST = BLOCK[O, CCB$A_FIRST_CIB; , BYTE],
CCB$S_HASH_BASE = BLOCK[O, CCB$V_HASH_BASE; , BYTE],
CCB$K_SPB_LIST = BLOCK[O, CCB$A_FIRST_SPB; , BYTE];
STRUCTURE
         CCB$HASH[I, O, P, S, E] =
```

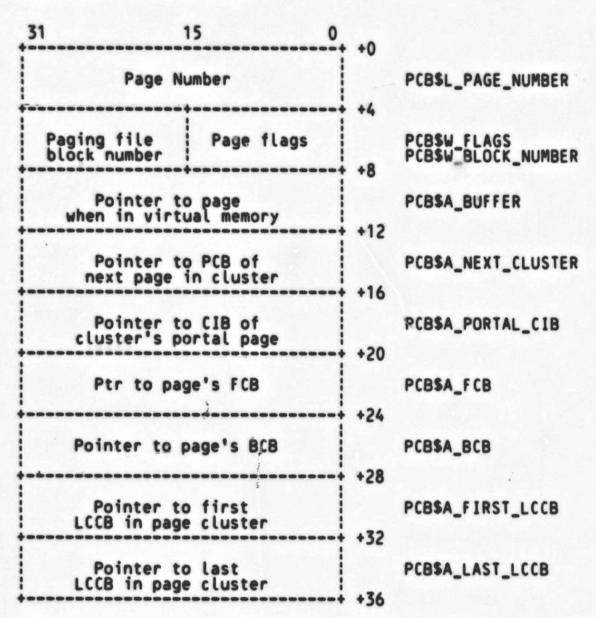
CDDLIB.B32;1

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(CCB\$HASH+CCB\$S\_HASH\_BASE+O+(I\*HCE\$S\_BLOCK\_LENGTH))<P,S,E>;

LITERAL CCBSM\_CORRUPT = 1^1 - 1^0; ! Stream is corrupt

Page Control Block (PCB)



A Page Control Block (PCB) exists for every page in the cache, and for pages that only have presence locks on them (thus they're not in the staging cache). Portal pages have a Cluster Information Block (CIB) appending to their PCB.

PCB\$S\_BLOCK\_LENGTH = 36;

MACRO
SPCB = BLOCK[PCB\$S\_BLOCK\_LENGTH,BYTE] FIELD (PCB\$Z\_FIELDS)

1

```
FIELD PCB$Z_FIELDS =
                       PCB$L_PAGE_NUMBER
PCB$W_FLAGS
PCB$V_MODIFIED
PCB$V_NEW_PAGE
PCB$V_READ_ONLY
PCB$V_FREE_PAGE
PCB$V_AVAILABLE
PCB$V_PORTAL_PAGE
PCB$W_BLOCK_NUMBER
PCB$A_BUFFER
PCB$A_NEXT_CLUSTER
PCB$A_PORTAL_CIB
PCB$A_FCB
PCB$A_FCB
PCB$A_FCB
PCB$A_FCB
PCB$A_LAST_CCCB
                                                                                                  = [0. 0. 32. 0],

= [4. 0. 16. 0],

= [4. 0. 1. 0],

= [4. 3. 1. 0],

= [4. 4. 1. 0],

= [4. 5. 1. 0],

= [4. 6. 1. 0],

= [4. 7. 1. 0],

= [6. 0. 16. 0],

= [8. 0. 32. 0],

= [12. 0. 32. 0],

= [120. 0. 32. 0],

= [24. 0. 32. 0],

= [24. 0. 32. 0],

= [28. 0. 32. 0],

= [32. 0. 32. 0],

= [32. 0. 32. 0],
                                                                                                                                                                      Must write back to dict.
Page was in free chain
                                                                                                                                                                      Page is read only
                                                                                                                                                               ! Page is a free page
! Page available in cache
! Page is portal page, CIB follows
             TES:
LITERAL
            PCB$K_LCCB_LIST
                                                                                                    = BLOCK[O, PCB$A_FIRST_LCCB; , BYTE];
LITERAL
           PCB$M_MODIFIED
PCB$M_NEW_PAGE
PCB$M_READ_ONLY
PCB$M_FREE_PAGE
PCB$M_AVAILABLE
PCB$M_PORTAL_PAGE
                                                                                                  = 1^1 - 1^0,
= 1^4 - 1^3,
= 1^5 - 1^4,
= 1^6 - 1^5,
= 1^7 - 1^6,
= 1^8 - 1^7;
                                                                                                                                                            Must write back to dict.
                                                                                                                                                             Page was in free chain
                                                                                                                                                             Page is read only
                                                                                                                                                           Page is a free page
Page available in cache
                                                                                                                                                           Portal page, CIB follows
```

# Cluster Information Block (CIB)

31	15 0	+ +0	
	Cluster flags	+4	CIB\$W_FLAGS CIB\$V_REF_COUNTS
Retrieval lock ref count	Presence lock ref count	+8	CIBSW_PRESENCE_COUNT CIBSW_RETRIEVAL_COUNT
Delete lock ref count	Update lock ref count	+12	CIBSW_UPDATE_COUNT CIBSW_DELETE_COUNT
Pointer to first deleted child cluster			CIBSA_FIRST_DELETED_CIB
Pointer to last deleted child cluster		+16	CIB\$A_LAST_DELETED_CIB
Ptr to newest lock granted for this page		+ +20	CIB\$A_LAST_LCB
Ptr to oldest lock granted for this page		+24	CIB\$A_FIRST_LCB
Ptr to next cluster at this directory level		+28	CIB\$A_NEXT_SIBLING
Ptr to prior cluster at this directory level		+32	CIB\$A_PRIOR_SIBLING
Ptr to first cluster		+36	CIB\$A_FIRST_CHILD
Ptr to last cluster at next directory level		+40	CIB\$A_LAST_CHILD
Ptr to CIB in next		+44	CIB\$A_PARENT_CIB
Ptr to portal CIB		+48	CIB\$A_HISTORY_CIB

LITERAL

%;

SET

FIELD

```
lin cluster holding history list !
            Ptr to NCB for the
                                             CIB$A_OWNER_NCB
       node that owns this cluster
       Ptr to cluster's pool header
                                             CIB$A_POOL
    Each clusters' portal page has a CIB associated with it. This block is physically appended to the PCB of the cluster's
    portal page.
    The CIB also points to the cluster's pool header.
    The CIB$V_REF_COUNTS lists must have each of the lock ref counts
    in the same order as the LOCK$K_xxx lock request symbols.
CIB$S_BLOCK_LENGTH = 60 + PCB$S_BLOCK_LENGTH;
$CIB = BLOCK[CIB$S_BLOCK_LENGTH,BYTE] FIELD (CIB$Z_FIELDS, PCB$Z_FIELDS)
   CIB$Z_FIELDS =
```

!\*

#### Lock Control Block (LCB)

LCBs are used to keep track of which locks exist on a cluster (CIB).

Each LCB is linked to its CIB, and to other LCBs for that cluster. When an LCB is granted, it is placed in the transaction's lock list until the transaction terminates.

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LITERAL

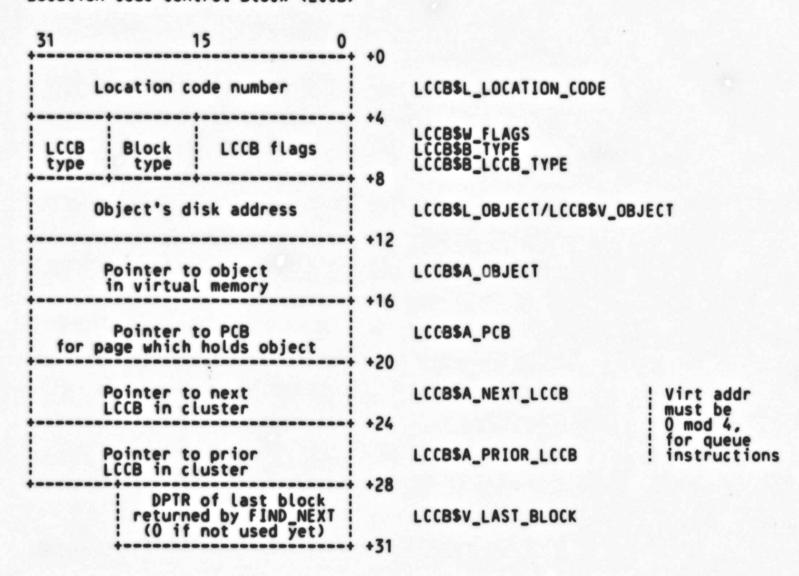
LCB\$K\_LOCK\_LIST = BLOCK[O, LCB\$A\_PRIOR\_LCB; , BYTE];

LCB\$K\_TEMP\_LOCK\_LIST = BLOCK[O, LCB\$A\_NEXT\_[OCK; , BYTE];

LITERAL LCBSM\_CURRENT

= 1^2 - 1^1; ! Allocated in current transaction

Location Code Control Block (LCCB)



Location codes are used to provide a convenient means for the user to identify a particular object. Each location code is associated with an LCCB. The following objects may be assigned location code:

- 1) Entities
- 2) Lists
- 3) Nodes

LCCB\$S\_BLOCK\_LENGTH = 31;

MACRO

```
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CDDLIB.B32:1
          $LCCB = BLOCK[LCCB$S_BLOCK_LENGTH,BYTE] FIELD (LCCB$Z_FIELDS)
         SET LCCBSZ_FIELDS =
FIELD
                  = [0. 0. 32. 0].
= [4. 0. 16. 0].
= [4. 0. 1. 0].
= [4. 1. 1. 0].
= [4. 3. 1. 0].
= [4. 3. 1. 0].
= [4. 5. 1. 0].
= [4. 6. 1. 0].
= [4. 8. 1. 0].
= [4. 8. 1. 0].
= [4. 8. 1. 0].
= [6. 0. 8. 0].
= [7. 0. 8. 0].
= [8. 0. 32. 0].
= [12. 0. 32. 0].
= [12. 0. 32. 0].
= [24. 0. 32. 0].
= [28. 0. 24. 0].
          TES:
LITERAL
         LCCB$K_LCCB_LIST
                                                                               = BLOCK[O, LCCB$A_NEXT_LCCB; , BYTE];
LITERAL
        LCCBSM_AVAILABLE
LCCBSM_MUST_SCAN
LCCBSM_GHOST
LCCBSM_DIRECTORY
LCCBSM_TERMINAL
LCCBSM_ENTITY_ATT
LCCBSM_ENTITY_LIST_ATT
LCCBSM_ENTITY_LIST_ATT
LCCBSM_ENTITY_LIST_ATT
LCCBSM_STRING_LIST_ATT
                                                                                                                            Object's virtual addr is known
                                                                             = 1°1 - 1°0,
= 1°2 - 1°1,
= 1°3 - 1°2,
= 1°4 - 1°3,
= 1°5 - 1°4,
= 1°6 - 1°5,
= 1°7 - 1°6,
= 1°8 - 1°7,
= 1°9 - 1°8,
                                                                                                                            Protection tree must be scanned
                                                                                                                           LCCB may not be fetched by LCC
Directory NCB
Terminal NCB
                                                                                                                       Entity attribute LCCB
Entity list attribute LCCB
Entity list ECCB
String list LCCB
         LCCB$K_LCCB_TYPE_FIRST
LCCB$K_NCB
LCCB$K_ECCB
LCCB$K_LCCB
LCCB$K_LCCB
                                                                              = 1,
                                                                                                                      ! LCCB includes NCB
! LCCB includes ECCB
! LCCB
```

```
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CDDLIB.B32:1
          Entity Cell Control Block (ECCB)
           31
                                15
                                    Cell number
                                                             ECCB$W_CELL
                         DPTR of this cell's
LST block
                                                             ECCB$V_LST
                    Ptr to LCCB of the
                                                              ECCB$A_PARENT_LCCB
                list that owns this cell
          Each cell in an entity list may be assigned a location code.
          This block is appended to the location code's LCCB to name the specific cell represented by the location code.
LITERAL
                                        LCCB$S_BLOCK_LENGTH + 9;
    ECCB$S_BLOCK_LENGTH =
MACRO
    SECCB = BLOCK[ECCB$S_BLOCK_LENGTH,BYTE] FIELD (ECCB$Z_FIELDS)
    %;
        ECCB$Z_FIELDS =
FIELD
     SET
         ECCB$W_CELL
ECCB$V_LST
ECCB$A_PARENT_LCCB
                                        = [0+LCCB$S_BLOCK_LENGTH, 0, 16, 0],
= [2+LCCB$S_BLOCK_LENGTH, 0, 24, 0],
= [5+LCCB$S_BLOCK_LENGTH, 0, 32, 0]
     TES:
```

### Node Control Block (NCB)

15	0 +0	
access rights this node		NCB\$L_ACCESS_RIGHTS
rights granted is node's ACL		NCB\$L_ACCESS_GRANTED
rights denied	*8	NCB\$L_ACCESS_DENIED
rights banished	+12	NCB\$L_ACCESS_BANISHED
is node's ACL	+16	
er to parent CB block		NCB\$A_PARENT_NCB
password		NCB\$V_PASSWORD NCB\$W_PSW_LENGTH NCB\$B_PSW_DTYPE NCB\$B_PSW_CLASS
	+24	NCB\$A_PSW_POINTER
e's   Entity name'	+ +28 s	NCB\$V_NAME NCB\$W_NAME_LENGTH NCB\$B_NAME_DTYPE
length	+32	NCB\$B_NAME_CLASS
er to entity's ame string	+ +36	NCB\$A_NAME_POINTER
		NCB\$A_SPB
	access rights this node  rights granted is node's ACL  rights denied is node's ACL  rights banished is node's ACL  er to parent CB block  rd's! Entity's a password e length  er to entity's sword string  e's Entity name' length  er to entity's ame string	access rights this node  rights granted is node's ACL  rights denied is node's ACL  rights banished is node's ACL  +12  rights banished is node's ACL  +16  er to parent CB block  rd's Entity's a password length  +24  er to entity's sword string  e's Entity name's length  +32  er to entity's ame string  +36

The LCCB for a node is followed by an NCB. This gives additional information needed for the node.

LITERAL NCB\$S\_BLOCK\_LENGTH = 40+LCCB\$S\_BLOCK\_LENGTH;

MACRO

```
CDDLIB.B32;1

$NCB = BLOCK[NCB$S_BLOCK_LENGTH,BYTE] FIELD (LCCB$Z_FIELDS, NCB$Z_FIELDS)

$;

FIELD NCB$Z_FIELDS =

NCB$L_ACCESS_RIGHTS = [0+LCCB$S_BLOCK_LENGTH, 0, 32, 0], NCB$L_ACCESS_GRANTED = [4+LCCB$S_BLOCK_LENGTH, 0, 32, 0], NCB$L_ACCESS_DENIED = [8+LCCB$S_BLOCK_LENGTH, 0, 32, 0], NCB$L_ACCESS_BANISHED = [12+LCCB$S_BLOCK_LENGTH, 0, 32, 0], NCB$L_ACCESS_BANISHED = [20+LCCB$S_BLOCK_LENGTH, 0, 32, 0], NCB$L_ACCESS_BANISHED = [22+LCCB$S_BLOCK_LENGTH, 0, 32, 0], NCB$L_ACCESS_BANISHED = [23+LCCB$S_BLOCK_LENGTH, 0, 32, 0], NCB$L_ACCESS_BANISH
```

### File Control Block (FCB)

31		15	0,	+0	
File	's flags	Channe numbe		+4	FCB\$W_CHAN FCB\$W_FLAGS
file Number			0	+8	FCB\$V_CDD\$TOP FCB\$B_FILE_NUMBER
	First 2 we	ords of		+12	FCB\$V_FID
	ID from manager	Last wo file's	rd of FID	+16	FCB\$W_LOCK_ID
	Pointer to	first free pag			FCB\$A_FREE_PAGE
Name's class	Name's data type		's i	+20	FCB\$V_FILENAME FCB\$W_FILE_LENGTH FCB\$B_FILE_DTYPE FCB\$B_FILE_CLASS
	Pointer to filename	entity's string		+28	FCB\$A_FILE_POINTER
	DKEY of bi	ock which			FCB\$L_OWNER
Name's class	Name's data type		me :	+32	FCB\$V_FULLNAME FCB\$W_FULL_LENGTH FCB\$B_FULL_DTYPE FCB\$B_FULL_CLASS
	Pointer to full filer	entity's	9	+40	FCB\$A_FULL_POINTER
P	ointer to free pag	temporary ge list		+44	FCB\$A_TEMP_FREE

Each active (open) dictionary file has an FCB.

The FCB\$L\_OWNER field holds the DKEY of the attribute block (FIL) which pointed to the file. The primary dictionary file has key 0, while files that are not currently pointed to have a value of -1

```
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  CDDLIB.B32:1
                            in the owner field.
FCB$K_FILE_LIMIT = 255,
FCB$S_BLOCK_LENGTH = 44;
  MACRO
              $FCB = BLOCK[FCB$S_BLOCK_LENGTH,BYTE] FIELD (FCB$Z_FIELDS)
 FIELD FCB$Z_FIELDS =
                          FCBSW_CHAN
FCBSW_FLAGS
FCBSV_READ_ONLY
FCBSV_ROOT
FCBSW_CDDSTOP
FCBSW_FID
FCBSW_FID
FCBSW_FID
FCBSW_FILE_NUMBER
FCBSW_FILE_NUMBER
FCBSW_FILE_NUMBER
FCBSW_FILE_NOME
FCBSW_FILE_LENGTH
FCBSW_FILE_CLASS
FCBSA_FILE_POINTER
FCBSW_FULL_LENGTH
FCBSW_FULL_LENGTH
FCBSW_FULL_LENGTH
FCBSW_FULL_LENGTH
FCBSW_FULL_LENGTH
FCBSB_FULL_CLASS
FCBSA_FULL_DTYPE
FCBSB_FULL_CLASS
FCBSA_FULL_POINTER
FCBSA_TEMP_FREE
                                                                                                       = [0, 0, 16, 0],

= [2, 0, 16, 0],

= [2, 0, 1, 0],

= [2, 1, 1, 0],

= [4, 0, 24, 0],

= [7, 0, 8, 0],

= [8, 0, 0, 0],

= [14, 0, 32, 0],

= [16, 0, 32, 0],

= [20, 0, 16, 0],

= [22, 0, 8, 0],

= [22, 0, 8, 0],

= [23, 0, 8, 0],

= [32, 0, 16, 0],

= [32, 0, 16, 0],

= [32, 0, 16, 0],

= [32, 0, 16, 0],

= [34, 0, 8, 0],

= [35, 0, 8, 0],

= [36, 0, 32, 0],

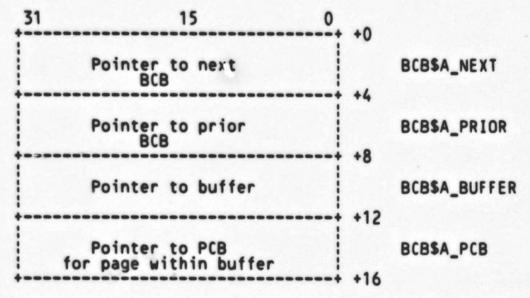
= [40, 0, 32, 0],
              TES:
 LITERAL
              FCB$M_READ_ONLY
                                                                              = 1^1 - 1^0,
                                                                                                                                   ! File can only be opened for read
                                                                                                                                  ! FCB is the root dictionary file
              FCB$M_ROOT
```

```
16-SEP-1984 16:58:51.80 Page 54
 CDDLIB.B32:1
           Pre-Allocated Page Block (PAPB)
            31
                      Pointer to next
                                                                 PAPB$A_NEXT
                      Number of free page
                                                                 PAPB$L_NUMBER
           Pre-allocated pages are represented by PAPBs linked onto the FCB in which the pages reside. Each PAPB has the page number of the free page it represents.
LITERAL PAPB$S_BLOCK_LENGTH = 8;
MACRO

$PAPB = BLOC PAPB$S_BLOCK_LENGTH, BYTE] FIELD (PAPB$Z_FIELDS)

%;
FIELD PAPB$Z_FIELDS =
          PAPB$A_NEXT
PAPB$L_NUMBER
                                                  = [0, 0, 32, 0],
= [4, 0, 32, 0]
     TES:
```

Buffer Control Block (BCB)



The buffer control blocks (BCBs) are used to control the pages that are in memory.

The BCBs are linked into a queue. When a buffer is needed, the last buffer in the queue is assigned.

The associated PCB contains information about the page and the buffer. If a page is modified while in the buffer, the page is written back to the work file before the buffer is reused.

```
LITERAL

BCB$K_NUMBER

BCB$S_BLOCK_LENGTH = 16;

MACRO

$BCB = BLOCK[BCB$S_BLOCK_LENGTH,BYTE] FIELD (BCB$Z_FIELDS)

%;

FIELD BCB$Z_FIELDS =

SET

BCB$A_NEXT

BCB$A_PRIOR

BCB$A_PRIOR

BCB$A_BUFFER

BCB$A_BUFFER

E [0, 0, 32, 0],

E [4, 0, 32, 0],

E [4, 0, 32, 0],

E [8, 0, 32, 0]
```

```
16-SEP-1984 16:58:51.80 Page 56
CDDLIB.B32;1
            User Retrieval List Entry Block
                                                    URLC$B_COUNT
                                        Count
                                                    URLC$V_UNUSED
           User Retrieval List Entry Block 7 0
                    Attribute Name
                                                    URLE$L_ATT/URLE$L_CELL
                      Cell Number
                                                    URLESW_DATA_TYPE
                                 Expected
                                                     URLE$W_RESERVE
                    Returned Status
                                                     URLE$L_STATUS
                                               + +12
URLE$W_USER_LENGTH
                                               URLESW_RETURN_LENGTH
                                User Buffer
                                                    URLE$A_USER_BUFFER
                    User Buffer or
                   String Descriptor
LITERAL
    URLC$S_BLOCK_LENGTH = 4,
URLE$S_BLOCK_LENGTH = 20;
    SURLC = BLOCK[URLC$S_BLOCK_LENGTH,BYTE] FIELD (URLC$Z_FIELDS)
FIELD
      URLC$Z_FIELDS =
        URLC$B_COUNT = [0, 0, 8, 0],
URLC$V_UNUSED = [1, 0, 24, 0]
    TES:
    $URLE = BLOCK[URLESS_BLOCK_LENGTH, BYTE] FIELD (URLESZ_FIELDS)
FIELD URLESZ_FIELDS =
```

```
16-SEP-1984 16:58:51.80 Page 58
 CDDLIB.B32:1
 .
            Write Control Block (WCB)
              31
                                       Number of pages
                                                                         WCB$W_SIZE
WCB$W_FLAGS
                      Flags
                          Ptr to page's permanant buffer
                                                                         WCB$A_OLD_BUF
                    Ptr to page's
location in write buffer
                                                                         WCB$A_NEW_BUF
                                        +12 WCB$V_IO_STATUS
               2nd word of 1st word of 1/0 status blk
                                                                         WCB$W_IO_STATUS
               4th word of 3rd word of 1/0 status blk
-
LITERAL WCB$S_BLOCK_LENGTH = 20;
      SWCB = BLOCK[WCB$S_BLOCK_LENGTH, BYTE] FIELD (WCB$Z_FIELDS)
FIELD
           WCB$Z_FIELDS =
      SET
           WCB$W_SIZE
WCB$W_FLAGS
WCB$V_WRITABLE
WCB$V_GROUP
WCB$A_OLD_BUF
WCB$A_NEW_BUF
WCB$V_IO_STATUS
WCB$W_IO_STATUS
                                                = [0, 0, 16, 0],

= [2, 0, 16, 0],

= [2, 0, 1, 0], ! Write group to file

= [2, 1, 1, 0], ! Group of pages

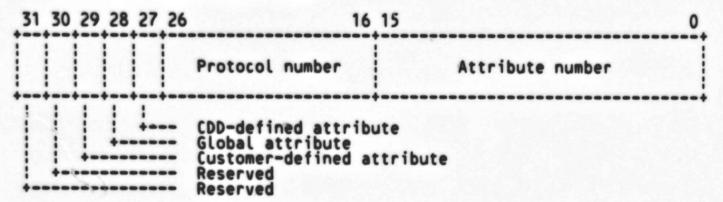
= [4, 0, 32, 0],

= [8, 0, 32, 0],

= [12, 0, 0, 0],

= [12, 0, 16, 0]
      TES:
LITERAL WCBSM_WRITABLE
                                    = 1^1 - 1^0;
                                                           ! Write group to file
! Group of pages
      WCB$M_GROUP
```

Attribute Name Block (ATNM)



The attribute name block defines the break down of the attribute name. Bits 0 to 15 contains the number. Bits 16 to 26 contains the protocol. The remaining bits are flags define the type of attribute. Bit 27 on implies a system defined attribute. Bit 28 on implies a global defined attribute. Bit 29 on implies a customer defined attribute.

NOTE: If the high order word of the attribute name block is 0 the block is describing a celi. That is the protocol will equal 0 and the flag bits will equal 0.

ATNM\$S\_BLOCK\_LENGTH = 4;
MACRO

SATNM = BLOCK[ATNM\$S\_BLOCK\_LENGTH,BYTE] FIELD (ATNM\$Z\_FIELDS)

FIELD ATNM\$Z\_FIELDS =

SET

ATNM\$W\_NUMBER = [0, 0, 16, 0],

ATNM\$W\_HIGH\_ORD = [0, 16, 16, 0],

ATNM\$V\_PROTOCOL = [0, 16, 11, 0],

ATNM\$V\_FLAGS = [0, 27, 5, 0],

ATNM\$V\_SYSTEM = [0, 27, 1, 0],

ATNM\$V\_GLOBAL = [0, 28, 1, 0],

ATNM\$V\_CUSTOMER = [0, 29, 1, 0]

```
16-SEP-1984 16:58:51.80 Page 60
CDDLIB.B32;1
           Page Number Block (PNB)
      31
                                               18 17
                Page Group Size
                                                                   Page Number
           The page number block define the break down of the page number. Bits 0 to 17 is the page number. Bits 18 to 31 define the number of consecutive pages that can be read or written at one time.
LITERAL
     PNB$S_BLOCK_LENGTH = 4;
MACRO
     $PNB = BLOCK[PNB$S_BLOCK_LENGTH, BYTE] FIELD (PNB$Z_FIELDS)
FIELD
         PNB$Z_FIELDS =
     SET
           PNB$V_PAGE_NUM = [0, 0, 17, 0]
PNB$V_GROUP_SIZ = [0, 17, 15, 0]
     TES:
```

SPB\$S\_BLOCK\_LENGTH = 37;

MACRO

!-

```
CDDLIB.B32;1

$SPB = BLOCK[SPB$S_BLOCK_LENGTH, BYTE] FIELD (SPB$Z_FIELDS)

$:

FIELD SPB$Z_FIELDS =

SET

$PB$A_NEXT_SPB = [0, 0, 32, 0],
$PB$A_PRIOR_SPB = [4, 0, 32, 0],
$PB$A_NCB = [8, 0, 32, 0],
$PB$L_GRANT = [12, 0, 32, 0],
$PB$L_DENY = [16, 0, 32, 0],
$PB$L_DENY = [16, 0, 32, 0],
$PB$L_ACCESS_RIGHTS = [24, 0, 32, 0],
$PB$V_PASSWORD = [28, 0, 0, 0],
$PB$W_PSW_LENGTH = [28, 0, 16, 0],
$PB$W_PSW_LENGTH = [28, 0, 16, 0],
$PB$B_PSW_CLASS = [31, 0, 8, 0],
$PB$B_PSW_CLASS = [31, 0, 8, 0],
$PB$B_PSW_POINTER = [32, 0, 32, 0],
$PB$B_SCAR_STATUS = [36, 0, 8, 0]

TES;

LITERAL
$PB$K_QUE_HEADER = BLOCK[0, SPB$A_NEXT_SPB; , BYTE];
```

```
16-SEP-1984 16:58:51.80 Page 63
CDDLIB.B32:1
        %SBTTL
                                'System Literal Definitions'
1++
                SYSTEM LITERAL DEFINITIONS
                These literals are only used internally.
!--
               CDD Implementation Version
LITERAL
       CDDSK_FACILITY
CDDSK_LOWEST_VERSION
CDDSK_VERSION
                                                              = 43
= 201;
= 201;
                                                                             ! Lowest compatible version ! Present version
               Boolean literals
LITERAL
       TRUE = 1.
FALSE = 0;
                                               ! Boolean TRUE value
                                               ! Boolean FALSE value
!+
                The following literals are used to validate routines' parameter
LITERAL
      ARGSK OPTIONAL
ARGSK REQUIRED
ARGSK SYNCIF
ARGSK MARK
ARGSK DEFAULT
ARGSK LONG
ARGSK WORD
ARGSK REF
ARGSK VALUE
                                                              Parameter is optional
Parameter is required
See CDD$$U_VALIDATE documentation
See CDD$$U_VALIDATE documentation
Mark parameter as missing
Use default value
Default value is a null string
Default value is a longword
Default value is a word value
Parameter passed by reference
Parameter passed by value
                                              = 12340.
               Access Lock Constants
                               The order of the LOCK$K_NORMAL lock constants MUST be the same as the order of the CIB$V_REF_COUNTS lock ref count fields in the CIB.
                NOTE:
-
```

```
16-SEP-1984 16:58:51.80 Page 64
CDDLIB.B32:1
LITERAL
      LOCKSK_NULL
                                                     = 0.
                                                                                ! Not locked
     LOCKSK_NORMAL
LOCKSK_PRESENCE
LOCKSK_RETRIEVAL
LOCKSK_UPDATE
LOCKSK_DELETE
LOCKSK_NORMAL_TYPES
                                                                                ! Base of partially queued locks
                                                     = 1.
                                                     = 4.
                                                     = 4.
                                                                                ! Number of partially queued locks
                                                     = 5.
           LOCK$K_QUEUED
LOCK$K_FULL_RET
LOCK$K_FULL_UPD
LOCK$K_QUEUED_TYPES
                                                                                ! Base of fully queued locks
! Number of fully queued locks
                                                     = 1^17-1^16,
= 1^18-1^17;
= 1^32-1^16;
      LOCKSK_XXX_GET_IF
LOCKSK_OPTIONS
                                                                                ! Flag indicates must establish lock ! Get lock if not already present ! Option bits for locks
1+
             Types of page purge requests
-
LITERAL
     PURGESK_ALL
PURGESK_PRESENCE
PURGESK_RETRIEVAL
PURGESK_UPDATE
PURGESK_SUBTREE
PURGESK_CLUSTER
                                       = 1.
= 2.
= 3.
= 4.
= 100.
                                                                   Complete purge
Purge but keep portal page
Checkpoint & keep retrieval locks
Checkpoint & keep all locks
                                                                      Purge whole subtree
                                        = 101:
                                                                   ! Purge this cluster only
             Types of blocks that can be allocated in a pool.
             Types of pools
LITERAL
      MEMSK_LOWEST_POOL = 1,

MEMSK_CCB_POOL = 1,

MEMSK_CIB_POOL = 2,

MEMSK_HIGREST_POOL = 2;
                                                                   Lowest pool type
Pool for CCB and HCBs
Pool for cluster blocks
                                                                   ! Highest pool type
             CCB Pool block types
LITERAL
      MEMSK_CCB_LOWEST
                                                                   ! Lowest block type in CCB pool ! CCB block
             MEMSK_CCB_CCB
```

```
16-SEP-1984 16:58:51.80 Page 65
CDDLIB.B32:1
     MEMSK_CCB_LOW_SLOT
MEMSK_CCB_HCB
MEMSK_CCB_HIGHEST
                                                         Lowest block type in free block list HCB block
                                                       ! Highest block type in CCB pool
.
           Cluster Pool block types
LITERAL
    MEMSK_CIB_LOWEST

MEMSK_CIB_CIB

MEMSK_CIB_LOW_SLOT

MEMSK_CIB_PCB

MEMSK_CIB_LCB

MEMSK_CIB_LCCB

MEMSK_CIB_LCCB

MEMSK_CIB_LCCB
                                                          Lowest block type in cluster pool
                                            122345
                                                          Lowest block type in free block list
                                                         PCB,
LCB,
                                                          ECCB.
     MEMSK_CIB_NCB
MEMSK_CIB_SPB
MEMSK_CIB_HIGREST
                                            = 6.
                                                          NCB,
                                                          SPB.
                                                        ! Highest block type in cluster pool
           These flags tell the deletion routine how it is to handle the
           following cases:
          DELSK_FAST
                                 says that pointers do not have to be cleaned up,
                                 as the block they reside in is going to be deleted.
           DELSK_PRESERVE indicates that a directory node is merely to be
                                 emptied, and that its cluster is not to be deleted.
          DELSK_SUBDICTIONARY
                                           says that sub-files are to have their contents
                                 deleted.
!-
LITERAL
     DELSK_FAST = 1^1 -
DELSK_PRESERVE = 1^2 -
DELSK_SUBDICTIONARY = 1^3 -
*
           User Identification Criteria
-
LITERAL
     CDD$K_ACL_LOWEST = 1,
CDD$K_ACL_PASSWORD = 1,
CDD$K_ACL_TERMINAL = 2,
CDD$K_ACL_UIC = 3,
CDD$K_ACL_USERNAME = 4,
CDD$K_ACL_HIGHEST = 4;
                                              PASSWORD
                                               TERMINAL name or class
                                              UIC
                                            ! USERNAME
```

```
16-SEP-1984 16:58:51.80 Page 66
 CDDLIB.B32:1
           %SBTTL
                                          'Security Masks'
 !++
                     SECURITY MASKS
!*
                    CDD security bits
LITERAL
         CDD$K_PROT_C
CDD$K_PROT_D
CDD$K_PROT_G
CDD$K_PROT_H
CDD$K_PROT_P
CDD$K_PROT_S
CDD$K_PROT_U
CDD$K_PROT_X
CDD$K_PROT_X
                                                                                                             CONTROL access
LOCAL DELETE access
GLOBAL DELETE access
HISTORY list entry creation access
                                                              = 1<sup>2</sup>
= 1<sup>3</sup>
= 1<sup>4</sup>
= 1<sup>6</sup>
= 1<sup>7</sup>
                                                                                                            PASS THRU access
SEE (read) access
UPDATE terminal node access
EXTEND directory node access
FORWARDing directory creation allowed
                                                              = 1^8
+
                    Macro-security values
                                                             = 1^9 - 1^0,
= CDD$K PROT D OR
CDD$K PROT G,
= CDD$K PROT F OR
CDD$K PROT C OR
CDD$K PROT D OR
CDD$K PROT G OR
CDD$K PROT H OR
CDD$K PROT U OR
CDD$K PROT U OR
CDD$K PROT X OR
CDD$K PROT F,
          CDD$K_PROT_ANY
CDD$K_PROT_DELETE
          CDD$K_PROT_EXTEND
          CDD$K_PROT_UPDATE
                    Other processor security bits
                    VAX-11 Datatrieve
          CDD$K_DTR_PROT_E
CDD$K_DTR_PROT_R
CDD$K_DTR_PROT_M
CDD$K_DTR_PROT_W
                                                              = 1^17 - 1^16;
= 1^18 - 1^17;
= 1^19 - 1^18;
= 1^20 - 1^19;
                                                                                                             EXTEND file
READ file
MODIFY file
WRITE file
```

```
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CDDLIB.B32:1
       %SBTTL
                               'User Literal Definitions'
               USER LITERAL DEFINITIONS
               These symbols are needed by users of the program interface.
               System Defined Attribute Names
LITERAL
       CDD$K_SYSNAM_FLAGS = 1^28 OR 1^27 OR 0^16; ! Global/System-defined/Protocol=0
LITERAL
      CDD$K_FIRST_SYSNAM = 1 OR CDD$K_SYSNAM_FLAGS,
CDD$K_FILE = 1 OR CDD$K_SYSNAM_FLAGS,
CDD$K_HISTORY = 2 OR CDD$K_SYSNAM_FLAGS,
CDD$K_NAME = 3 OR CDD$K_SYSNAM_FLAGS,
CDD$K_PROTOCOL = 5 OR CDD$K_SYSNAM_FLAGS,
CDD$K_TYPE = 6 OR CDD$K_SYSNAM_FLAGS,
CDD$K_PATHNAME = 7 OR CDD$K_SYSNAM_FLAGS,
CDD$K_SHORT_PATHNAME = 8 OR CDD$K_SYSNAM_FLAGS,
CDD$K_ORDER = 9 OR CDD$K_SYSNAM_FLAGS,
CDD$K_LAST_SYSNAM = 9 OR CDD$K_SYSNAM_FLAGS,
                                                                                                               Lowest system defined attribute name value Node's file name
                                                                                                               History list head
                                                                                                               Node's name
Node's protocol name
Type of object pointed to by location code
                                                                                                               Node's complete pathname
Node's path to CDD$DEFAULT directory
                                                                                                               Directory's order
                                                                                                            ! Highest system defined attribute name value
               Attribute and Entity Types
LITERAL
      CDD$K_FIRST_TYPE
CDD$K_ENTITY
CDD$K_ENTITY_LIST
CDD$K_NUMERIC
CDD$K_STRING
CDD$K_STRING_LIST
CDD$K_DIRECTORY
CDD$K_TERMINAL
CDD$K_LAST_TYPE
                                              = 1.
= 2.
= 3.
                                              = 4.
                                              = 6.
= 7.
= 8.
= 8:
               User's entity purge options
LITERAL
       CDD$K_ALL
CDD$K_ABORT
CDD$K_CHECKPOINT
                                              = 1^2 - 1^1
```

```
CDDLIB.B32;1

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LITERAL
CDD$K_NOHISTORY = 1^1 - 1^0, | Doesn't want history list cluster CDD$K_NOACL = 1^2 - 1^1, | Don't create default ACL entry CDD$K_CREATE = 1^3 - 1^2, | Create dictionary file if needed CDD$K_FIRST = 1^4 - 1^3, | Insert as first node CDD$K_LAST = 1^5 - 1^4; | Insert as last node

LITERAL
CDD$K_SUBDICTIONARY = 1^2 - 1^1; | Pail if directory has children CDD$K_SUBDICTIONARY = 1^2 - 1^1; | Delete contents of subdictionaries

Values of the CDD$K_ORDER attribute

LITERAL
CDD$K_SORTED = 1, | Directory is sorted CDD$K_NONSORTED = 2; | Directory is not sorted
```

```
CDDLIB.B32:1

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**SBTTL 'LINKAGE DEFINITIONS'

LINKAGE DEFINITIONS

CDDCALL

This Linkage uses the CALLG/CALLS linkage convention, except that it allows for one global register to be used in parameter passing.

R11 -- used to pass the user's context pointer.

LINKAGE
CDDCALL = CALL : GLOBAL (USER_CONTEXT = 11);

**

SYS_JSB

This Linkage provides us with a general JSB routine linkage.

LINKAGE
SYS_JSB = JSB;
```

```
16-SEP-1984 16:58:51.80 Page 70
 CDDLIB.B32:1
       %SBTTL
                         "MACRO DEFINITIONS"
             MACRO DEFINITIONS
             SACTIVE
SINACTIVE
             These macroes declare that we have started, and finished, respectively, a CDD transaction. They abort the transaction if another transaction is in progress.
 !-
MACRO
SACTIVE =
BEGIN
                  EXTERNAL
                        CDD$GB_INUSE:
                                                           BYTE:
                  EXTERNAL LITERAL CDD$_NOTASTREE;
                  BUILTIN
                        TESTBITSS:
                  IF TESTBITSS (CDD$GB_INUSE) THEN SIGNAL (CDD$_NOTASTREE);
            END
       %.
       $INACTIVE =
             BEGIN
                  EXTERNAL
                        CDD$GB_INUSE:
                                                           BYTE:
                  CDD$GB_INUSE = FALSE;
            END
       Z.
             SBITCLEAR
             This macro checks to see if any bit in a mask is set in the target area. If not, it returns TRUE.
       $BITCLEAR(target, mask) = (target AND mask) EQLU 0
```

```
$BITSET
```

This macro checks to see if any bit in a mask is set in the target area. If so, it returns TRUE.

\$BITSET(target, mask) = (target AND mask) NEQU 0

## SDONE\_TRANS

This macro is used to terminate a transaction.

Call:

\$DONE\_TRANS [(dsc1 [, dsc2] ...)]

Where:

dsci :== the names of dynamic descriptors which are to have their strings returned to the string pool.

SDONE TRANS (dsc1) = BEGIN

EXTERNAL ROUTINE CDD\$\$SN\_DONE\_TRANS

: CDDCALL NOVALUE;

X.

!\*

## SFIND\_ENTITY

This macro returns the virtual address of the LCCB associated with a location code. It also checks to make certain that the cluster is locked as requested, and that the cluster's node allows the requested security access.

## Call:

ANY

Where:

valid-arg :== the address of the calling routine's validated
argument list.
The argument list must have the following format:

valid-arg[0] ::= address of longword holding
 context #
valid-arg[1] ::= address of descriptor holding
 path name, or zero.

The first set of keywords names the desired lock state of the entity's cluster.

The last set of keywords names the intended access to the cluster.

\$FIND\_ENTITY (valid\_arg, locking, security) =
BEGIN
EXTERNAL ROUTINE
CDD\$\$SN\_FIND\_ENTITY : CDDCALL;

CDD\$\$SN FIND ENTITY (...valid arg[1], \$FIND\_XXX\_LOCKING (%REMOVE(locking)), %NAME("CDD\$K\_PROT\_", security))

X. END

%.

.

\$FIND\_XXX\_LOCKING (class, type) = %IF %IDENTICAL (class, %QUOTE LOCK) %THEN LOCK\$K\_XXX\_GET OR

XELSE
XIF XIDENTICAL (class, XQUOTE LOCKIF) XTHEN
LOCKSK\_XXX\_GET\_IF OR
XELSE

XELSE
XIF NOT XIDENTICAL (class, %QUOTE CHECK) %THEN
XERROR ('Invalid locking keyword: ', class)

XFI XFI

%NAME ('LOCK\$K\_', type)

SFIND\_NODE

This macro returns the virtual address of the NCB associated with a path name or location code. It also checks to make certain that the cluster is locked as requested, and that the target node allows the requested security access.

```
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CDDLIB.B32:1
          Call:
              ncb-block.wa.v = $FIND_NODE (valid-arg.ra.v ,
LOCK RETRIEVAL READ
( { LOCKIF } { UPDATE } ) , { MODIFY DELETE
                                                       ) , { MODIFY
                                                                           )):
                                                                 ANY
         Where:
               valid-arg :== the address of the calling routine's validated
                                 argument list.
                                 The argument list must have the following format:
                                         valid-arg[0] ::= address of longword holding
                                                                context #
                                         valid-arg[1] ::= address of descriptor holding
                                         valid-arg[2] ::= address of longword holding
                                                                location code, or zero.
               The first set of keywords names the desired lock state of
               the node's cluster.
               The last set of keywords names the intended access to the
              cluster.
    $FIND_NODE (valid_arg, locking, security) =
         BEGIN
              EXTERNAL ROUTINE
                   CDD$$SN_FIND_NODE
                                                   : CDDCALL:
              CDD$$SN_FIND_NODE (valid_arg,
$FIND_XXX_LOCKING (%REMOVE(locking)),
%NAME('CDD$K_PROT_', security))
         END
    X.
         SFIND_PARENT
         This macro returns the virtual address of the NCB associated with a location code. It also checks to make certain that the cluster is locked as requested, and that the cluster's node allows the requested security access.
         Call:
               status.wlc.v = $FIND_PARENT (valid-arg.ra.v .
                         LOCK
                                         RETRIEVAL
                                                                 READ
                    ( { LOCKIF } ( UPDATE
                                                       )), { MODIFY }, ncb-block.wa.r,
                         CHECK
                                         DELETE
                                                                 DELETE
                                                                 ANY
                    name.wt.ds);
```

Where: valid-arg :== the address of the calling routine's validated
argument list. The argument list must have the following format: valid-arg[0] ::= address of longword holding context # valid-arg[1] ::= address of descriptor holding valid-arg[2] ::= address of longword holding location code, or zero. The first set of keywords names the desired lock state of the target cluster. The last set of keywords names the intended access to the cluster. \$FIND\_PARENT (valid\_arg, locking, security, ncb\_block, name) = BEGIN EXTERNAL ROUTINE CDD\$\$SN\_FIND\_PARENT : CDDCALL: CDD\$\$SN\_FIND\_PARENT (valid arg, \$FIND\_XXX\_LOCKING (%REMOVE(locking)), %NAME('CDD\$K\_PROT\_', security), ncb\_block, name) END SINIT\_DSC This macro is used to initialize dynamic string descriptors. Call: \$INIT\_DSC (dsc1 [, dsc2] ...) \$INIT\_DSC[DSC\_NAM] = BEGIN DSC\_NAM[DSC\$B\_DTYPE] = DSC\$K\_DTYPE\_T;
DSC\_NAM[DSC\$B\_CLASS] = DSC\$K\_CLASS\_D;
DSC\_NAM[DSC\$W\_LENGTH] = 0;
DSC\_NAM[DSC\$A\_POINTER] = 0; END

\$10\_SYNC (ef, iosb)

CDDLIB.B32:1

%.

Z.

!-

.

STATUS = \$WAITFR (efn = ef); IF NOT .STATUS THEN SIGNAL STOP (.STATUS); STATUS = SCLREF (efn = ef); IF NOT .STATUS THEN SIGNAL STOP (.STATUS); END WHILE .iosb[0] EQLU 0 %.

STATUS:

SMARK\_PAGE (page-block)

CDDLIB.B32:1

\*

+

ef

iosb

\$10\_SYNC (ef, iosb) = BEGIN

LOCAL

This macro marks a page as modified. Such a page must be written back to the dictionary file when it is purged from the staging buffers.

\$MARK\_PAGE (page\_block) =
 page\_block[PCB\$v\_MODIFIED] = TRUE

\$PARAMETERS (arg1 [, arg2] ...)

This macro is used to build the control vector for the parameter list validate routine (CDD\$\$U\_VALIDATE).

There is one entry (arg1, arg2, etc) for every formal parameter in the routine. Each entry has the following format:

( REQUIRED , [n] , [REF : VALUE] : ( SYNC : SYNCIF ) , n : OPTIONAL [,MARK : ,DEFAULT [,STRING : ,LONG : ,WORD] ] )

\$PARAMETERS[] =
 UPLIT (%LENGTH, \$PARM\_PROCESS (%REMOVE(%REMAINING)))

```
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CDDLIB.B32:1
     %.
     $PARM_PROCESS[arg] =
    BYTE ($PARM_DECIDE (%REMOVE(arg)))
    SPARM_DECIDE(status)[] =
%IF %IDENTICAL (status, %QUOTE REQUIRED) %THEN
ARG$K_REQUIRED, 0,
%IF %NULL (%REMAINING) %THEN
ARG$K_REF, 0
                      SPARM_REQUIRED (%REMAINING)
                XF I
          XELSE

XIF XIDENTICAL (status, XQUOTE SYNC) XTHEN

ARG$K_SYNC, 0, 0, XREMAINING

XELSE

XIDENTICAL (status, XQUOTE SYNCIF)
                      %IF %IDENTICAL (status, %QUOTE SYNCIF) %THEN ARG$K_SYNCIF, 0, 0, %REMAINING %ELSE
                           %IF %IDENTICAL (status, %QUOTE OPTIONAL) %THEN ARG$K_OPTIONAL, $PARM_OPTIONAL (%REMAINING), 0 %ELSE
                                 XERROR ('Invalid parameter status: ', status)
0, 0, 0
                            XF I
                     %FI
               XF I
          XF I
     $PARM_REQUIRED(number, pass) = 
%IF %NULL (pass) %THEN
           XELSE
                %NAME ('ARG$K_', pass)
           %IF %NULL (number) %THEN
          XELSÉ
                , number
          XF I
     %.
     $PARM_OPTIONAL(action)[] =
           XIF XIDENTICAL (action, XQUOTE MARK) XTHEN
                ARG$K_MARK, 0
           XELSE
                %IF %IDENTICAL (action, %QUOTE DEFAULT) %THEN
                      ARGSK_DEFAULT, %NAME ('ARGSK_', %REMAINING)
                XELSE
                      XERROR ('Invalid action keyword: ', action)
                      0. 0
                %FI
           XF I
```

SRECOVERY ( RESET ! ENABLE ! DISABLE )

This macro determines the ability of the exit handler to perform a purge of the cache if the task aborts.

SRECOVERY (DISABLE)

CDDLIB.B32:1

**SPRESENT** 

\$PRESENT (name) = . name NEQA 0

%.

!+

1-

1+

-

declares that the internal data structures or external disk structure is in an indeterminant state and cannot be recovered by the exit handler.

SRECOVERY (ENABLE)

declares that the data structure manipulation is completed.

**SRECOVERY (RESET)** 

specifies that the data structures are in a recoverable state.

Note that these can be nested. Recovery is only possible if the recovery counter is zero (reset).

\$RECOVERY (option) = BEGIN EXTERNAL CDD\$GW\_RECOVERY: WORD: %IF %IDENTICAL (option, %QUOTE DISABLE) %THEN CDD\$GW\_RECOVERY = .CDD\$GW\_RECOVERY + 1; %IF %IDENTICAL (option, %QUOTE ENABLE) %THEN CDD\$GW\_RECOVERY = .CDD\$GW\_RECOVERY - 1; %ELSE IF IDENTICAL (option, IQUOTE RESET) ITHEN CDDSGW\_RECOVERY = 0; **XERROR** ('Illegal recovery option: ', option) XF I XF I

```
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CDDLIB.B32;1
              %FI
         END
    %.
         SRELEASE_LOCK
         This macro calls the CDD$$SN_RELEASE routine to release one
         or more locks.
         Call:
             $RELEASE_LOCK ( { UPDATE } , ncb-block1 ...);
    *RELEASE LOCK (locking)[] = BEGIN
             EXTERNAL ROUTINE CDD$$SN_RELEASE
                                              : CDDCALL
                                                                NOVALUE:
             CDD$$SN_RELEASE (%NAME ('LOCK$K_', locking), %REMAINING)
         END
    %.
         $SIGNAL_SEVERE (error)
         This routine signals a severe error.
    $SIGNAL SEVERE (ERROR) = SIGNAL (ERROR OR STS$K_SEVERE)
.
         $STATIC_DSC
         This macro is used to initialize static string descriptors.
         Call:
             $STATIC_DSC (dsc1 [, dsc2] ...)
         Where:
             dsci :== name ! ( name , source )
         Source is the name of another descriptor. The named descriptor is initialized to point to the same string as source.
```

```
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CDDLIB.B32:1
!-
       $STATIC_DSC[dsci] = 
$STATIC_DSC_BUILD (%REMOVE (dsci))
       $STATIC_DSC_BUILD(name, source) = BEGIN
                    name[DSC$B_DTYPE] = DSC$K_DTYPE_T;
name[DSC$B_CLASS] = DSC$K_CLASS_S;
%IF %NULL(source) %THEN
    name[DSC$W_LENGTH] = 0;
    name[DSC$A_POINTER] = 0;
                            name[DSC$W_LENGTH] = .source[DSC$W_LENGTH];
name[DSC$A_POINTER] = .source[DSC$A_POINTER];
              END
       %.
!+
              $STRING
              These macros are used to build string descriptors for literal
              strings.
              Call:
                            $STRING ( (name, string) ...)
$STRING_INIT ();
             "name" is defined to be a BLOCK structure, and the address of
the string is poked into the structure by the STRING_INIT macro,
which must be the first executable statement in a routine.
-
      $STRING [] =
$STRING_DSC_SETUP (%REMAINING)
MACRO
                    *QUOTE *QUOTE *STRING_INIT = 
*STRING_PTR_SETUP (*QUOTE *EXPAND *REMAINING)
                     XQUOTE X
      X.
      $STRING_DSC_SETUP[PAIR] =
$STRING_DSC_INIT (%REMOVE(PAIR))
       $STRING_DSC_INIT (STR_NAME, STR_VAL) =
                    STR_NAME: $DSC PRESET ([DSC$B_DTYPE] = DSC$K_DTYPE_T,
[DSC$B_CLASS] = DSC$K_CLASS_S,
[DSC$W_LENGTH] =

%CHARCOUNT (%REMOVE(STR_VAL)));
```

```
CODLIB.B32:1
    %.
    $STRING_PTR_SETUP[PAIR] =
$STRING_PTR_INIT (%REMOVE(PAIR))
    STRING_PTR_INIT (STR_NAME, STR_VAL) = STR_NAME[DSC$A_POINTER] = UPLIT BYTE(%REMOVE(STR_VAL));
!+
         STEXTC
         This macro is used to define counted strings. The first byte
         of such strings is a count of the number of characters in the
         string.
         Each string is defined to be a VECTOR[,BYTE] structure, with the O element being the character count, and the actual string starting at STRING[1].
         $TEXTC ((name1,'str1') [, (name2, 'str2')] ...);
    STEXTC[PAIR] =
         STEXTC_STR(%REMOVE(PAIR))
    STEXTC STR(NAME, TSTR) = BIND
              NAME = UPLIT BYTE (%CHARCOUNT(%REMOVE(TSTR)), %REMOVE(TSTR)) :
                   VECTOR[%CHARCOUNT(%REMOVE(TSTR))+1, BYTE];
    %.
         $VALIDATE (cntl, [arg-list])
         This macro generates a call to the general transaction setup
         routine.
         cntl
                   -- the address of the control vector for CDD$$U_VALIDATE.
         arg-list-- is optional. If present, it is the address of the vector which is to receive the verified argument list
                       from CDD$$U_VALIDATE.
    $VALIDATE(cntl, arg_list) =
```

: CDDCALL NOVALUE:

BEGIN

EXTERNAL ROUTINE

CDD\$\$SN\_START\_TRANS

```
CDDLIB.B32;1

BUILTIN
AP;

XIF XNULL (arg_list) XTHEN
CDD$$SN_START_TRANS (.AP, cntl)
XELSE
CDD$$SN_START_TRANS (.AP, cntl, arg_list)

END
X;
```

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